

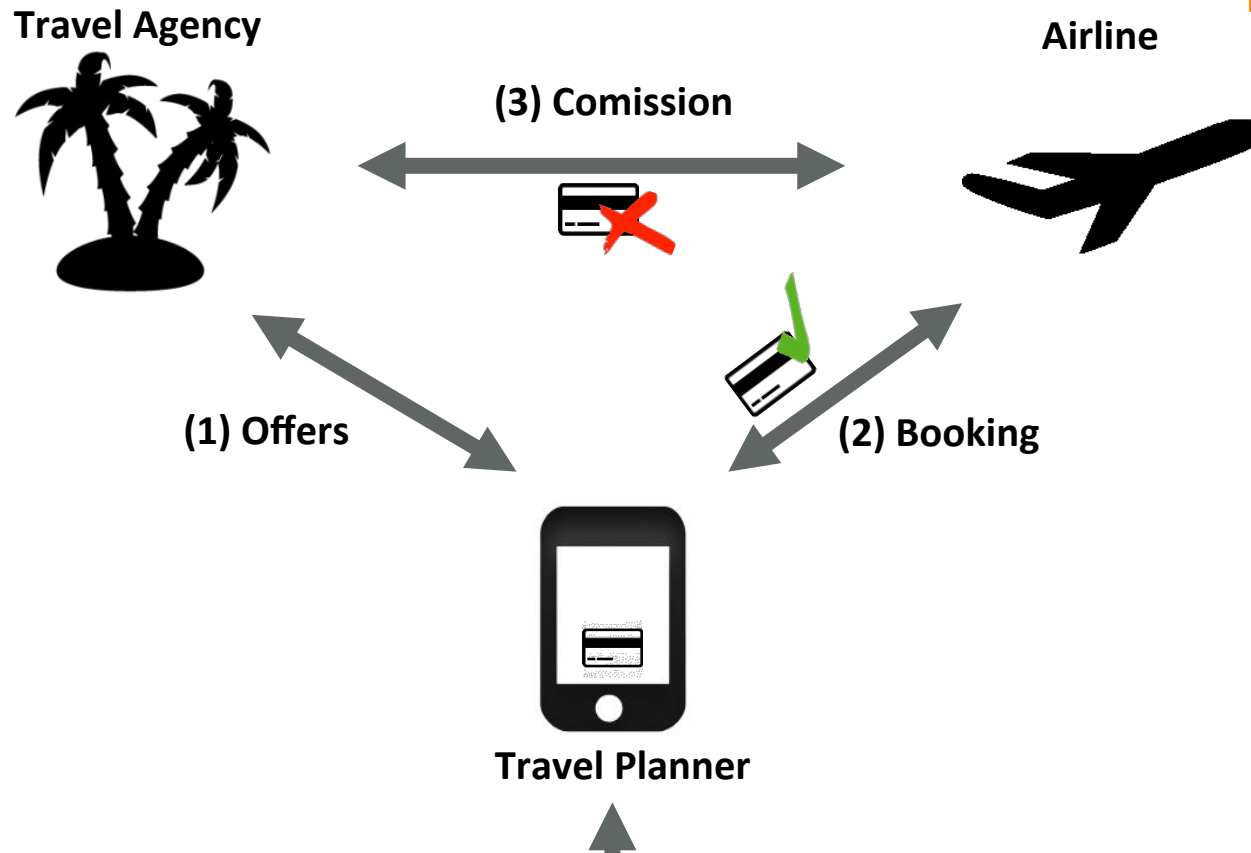


IFlow — Model driven development of information flow secure systems

Wolfgang Reif, Kuzman Katkalov, Marian Borek, Kurt Stenzel

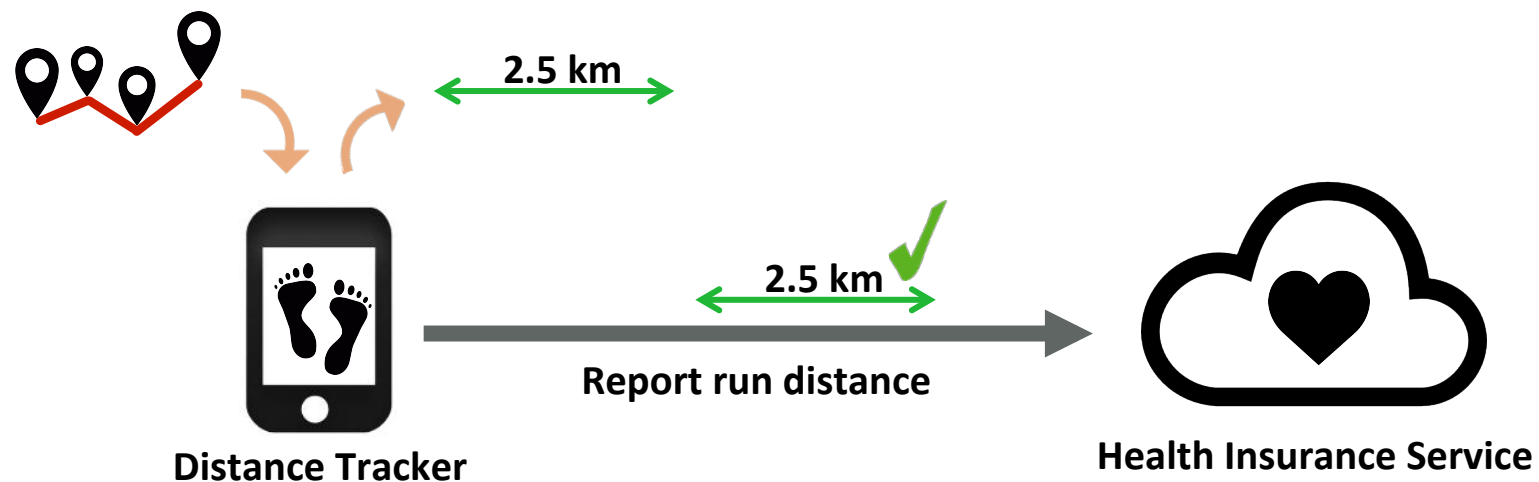


Where does your information go?



- Travel Agency **never learns** the user's credit card data
- Airline only learns the user's credit card data **after user confirmation**

How much information are you sharing?



- Health Insurance Service **never learns** the user's location
- Distance Tracker **properly anonymizes** the user's location data as distance

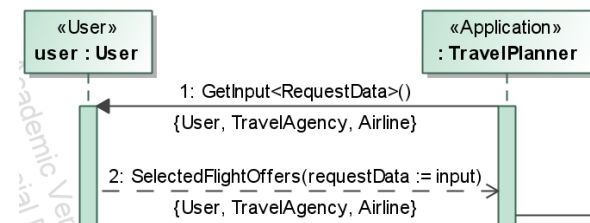
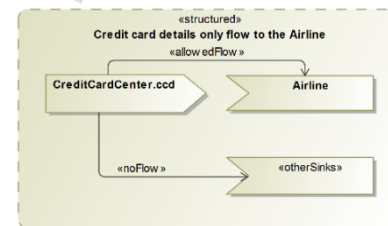
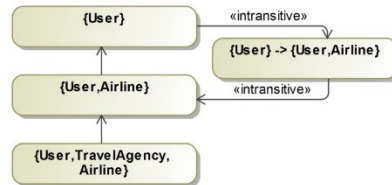
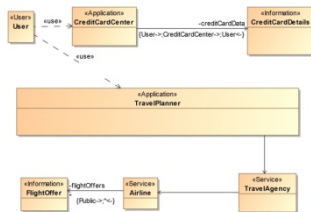
- **IFlow: a model driven approach for developing distributed apps with secure information flow**
- Software/security co-design in distributed apps
 - Modeling
 - Automatic check and verification
- Special aspect: limited release of information
 - Modeling
 - Verification

IFlow: a model-driven approach to IFC

MODEL > FLOW

Modeling language

- *principle: security by design*
- *graphical, domain specific language for IFC*
- *distributed apps & services*
- *with tool support*
- *and code generation*



Formal foundation

- *formal semantics*
(ASMs & predicate logic)
- *intransitive noninterference*
(extension of Rushby/Meyden)
- *declassification*
(“what”, “where”, “when”)



A hybrid approach to provide guarantees

Information flow:
Automatic code check

Travel Agency *never learns* the
user's credit card data

Health Insurance Service *never
learns* the user's location



Information release:
Formal verification

Airline only learns the user's
credit card data
after user confirmation

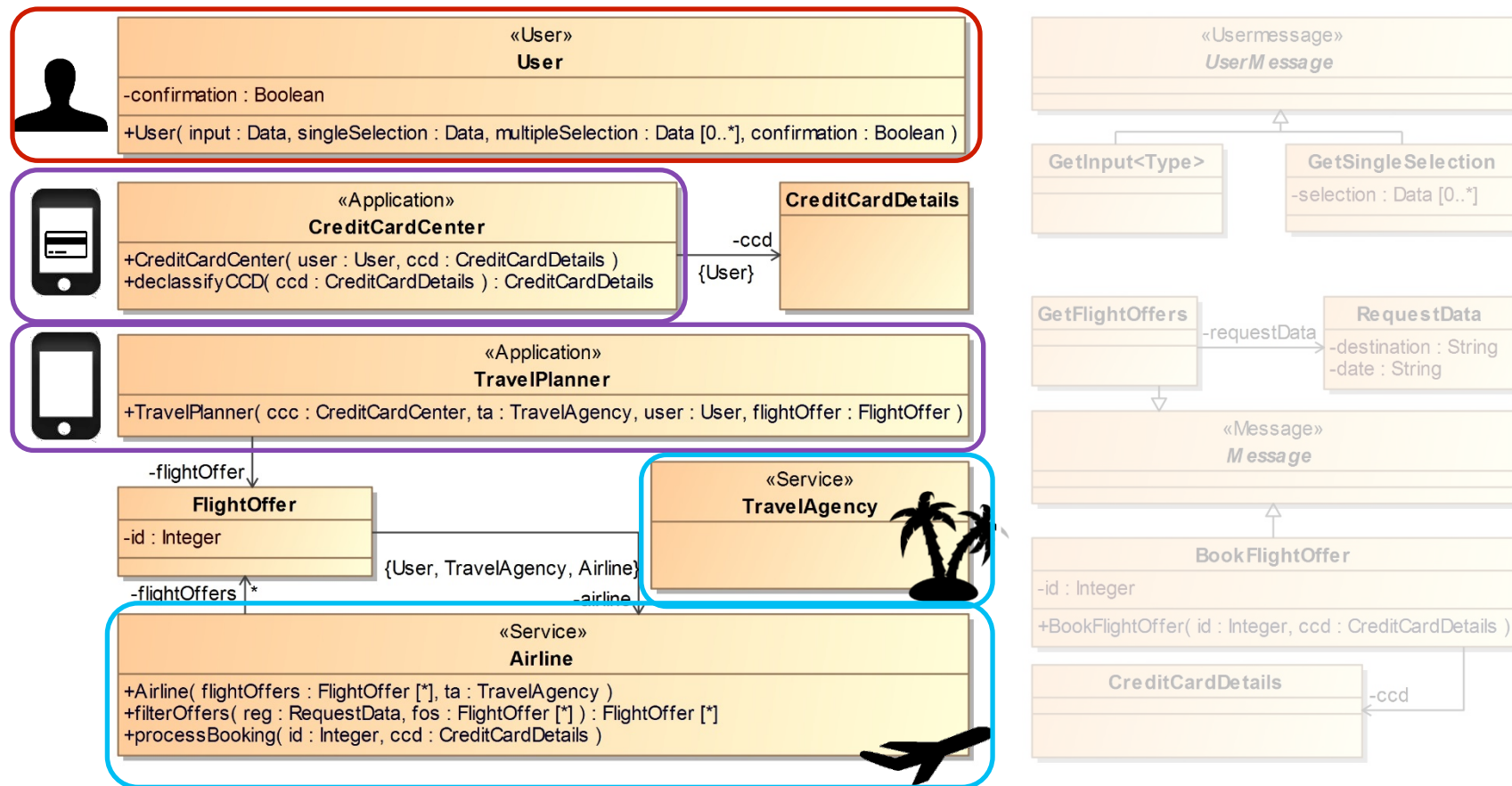
Distance Tracker *properly
anonymizes* the user's location



- IFlow: a model driven approach for developing distributed apps with secure information flow
- **Software/security co-design in distributed apps**
 - Modeling
 - Automatic check and verification
- Special aspect: limited release of information
 - Modeling
 - Verification

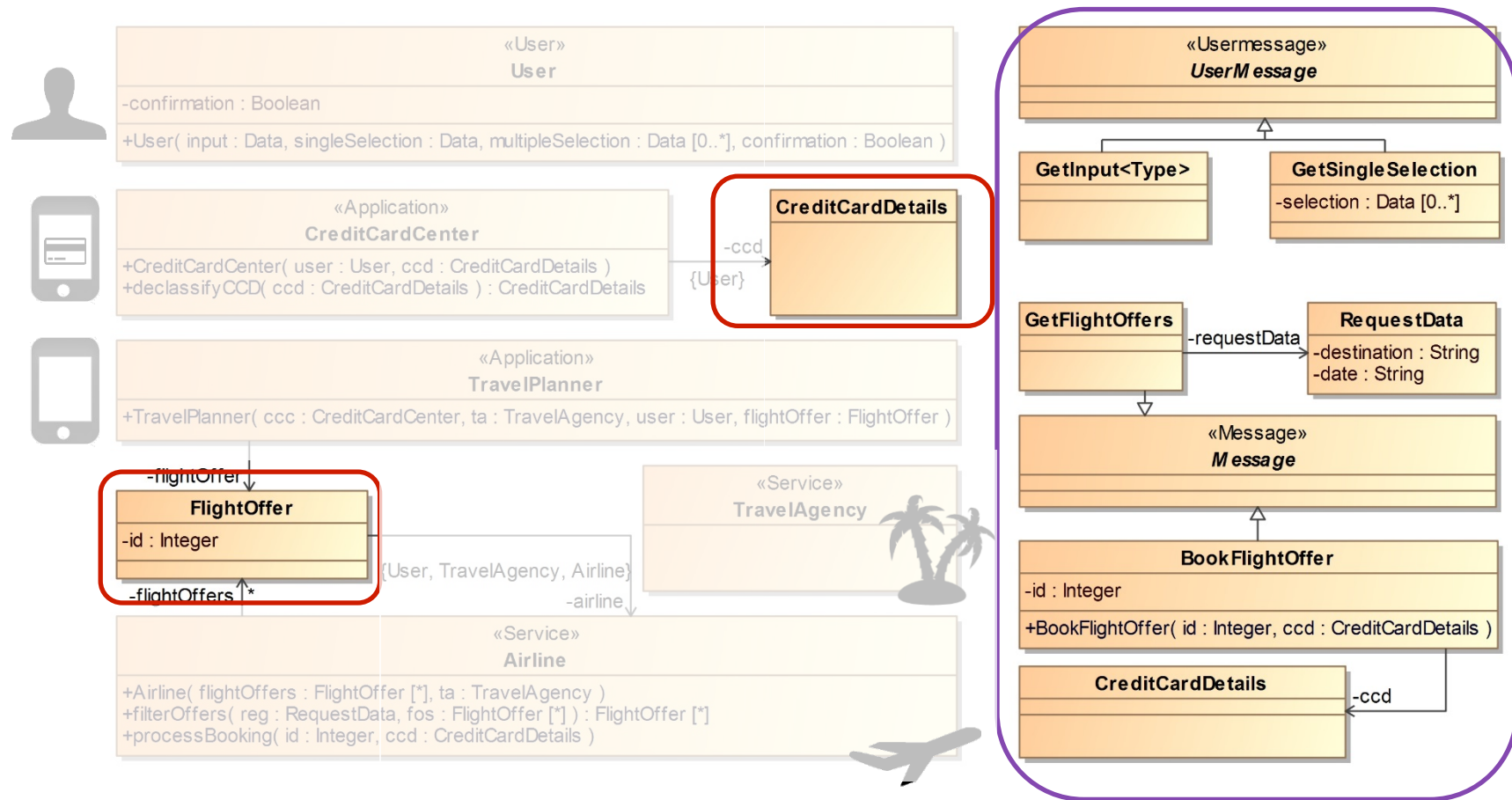
TravelPlanner: components

Application components: **user interface**, **mobile apps** and **web services**

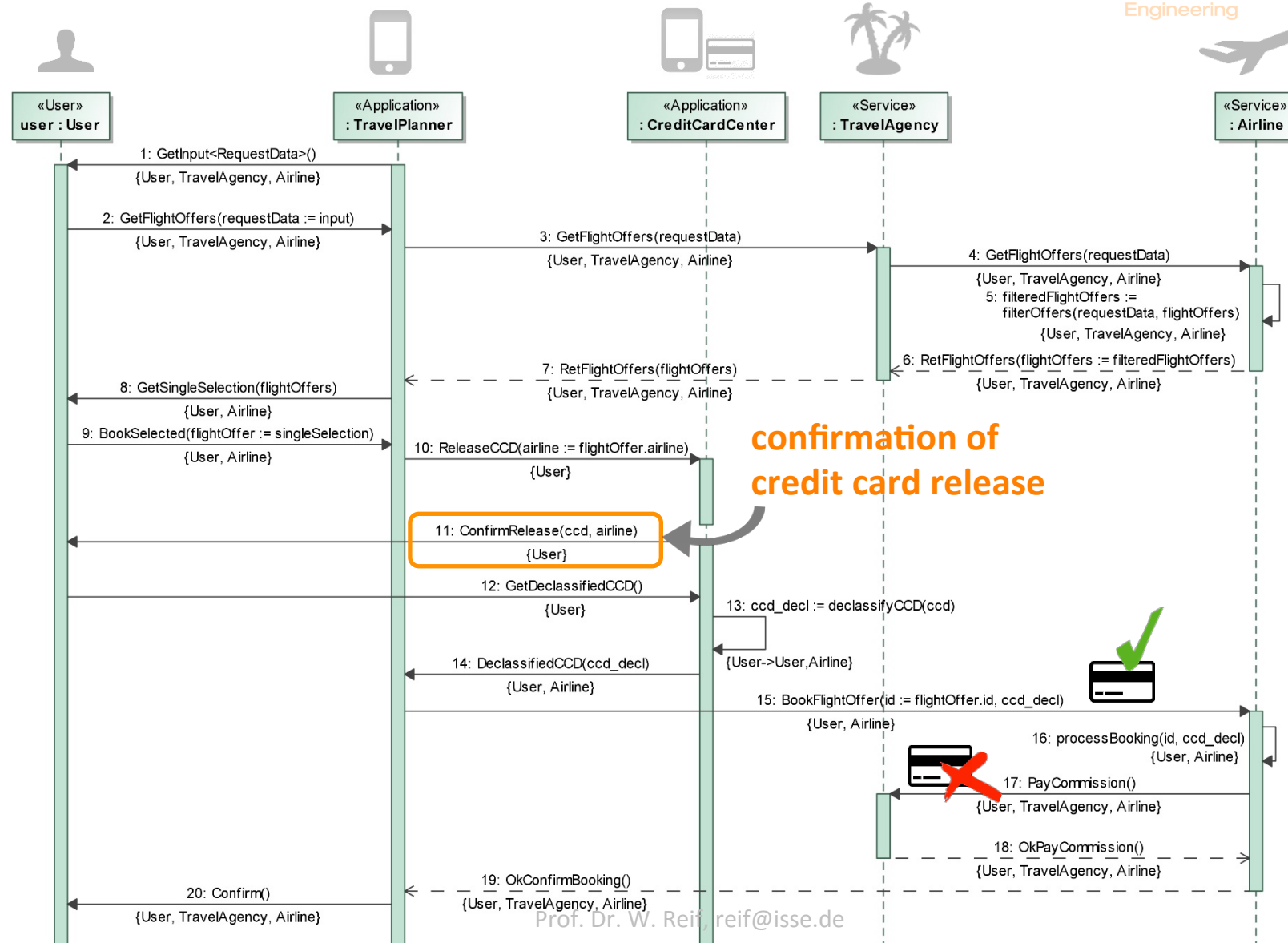


TravelPlanner: data types

Data types and message classes



TravelPlanner: behaviour



Travel Planner: security

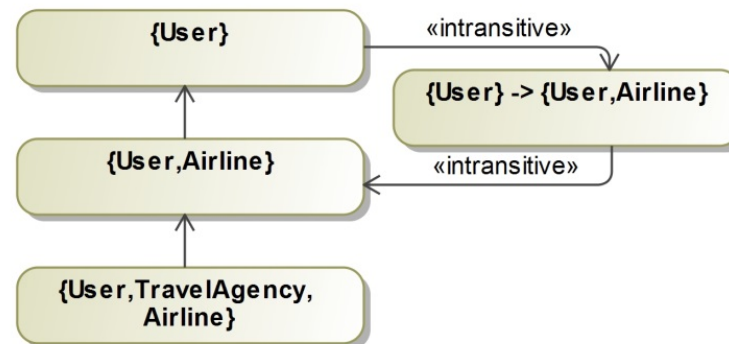
- **Subjects** (user, apps & web services): S

$$S_{\text{TravelPlanner}} = \{User, Airline, TravelAgency\}$$

- **Security domains:** $d \subseteq S, d \subseteq S \times S$

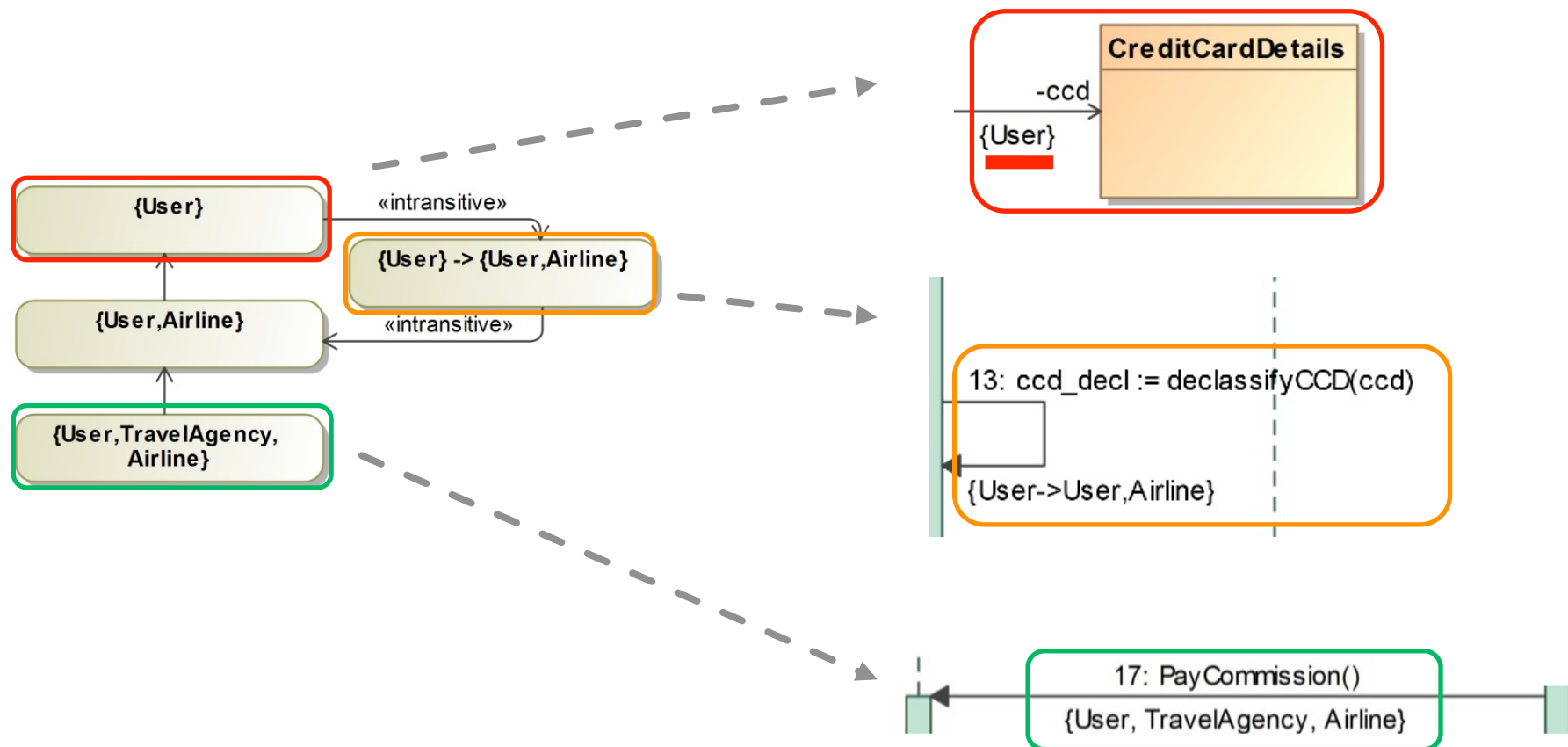
$\{User\}, \{User, Airline\}, \{User, Airline, TravelAgency\},$
 $\{User\} \rightarrow \{User, Airline\}$

- **Security policy:**

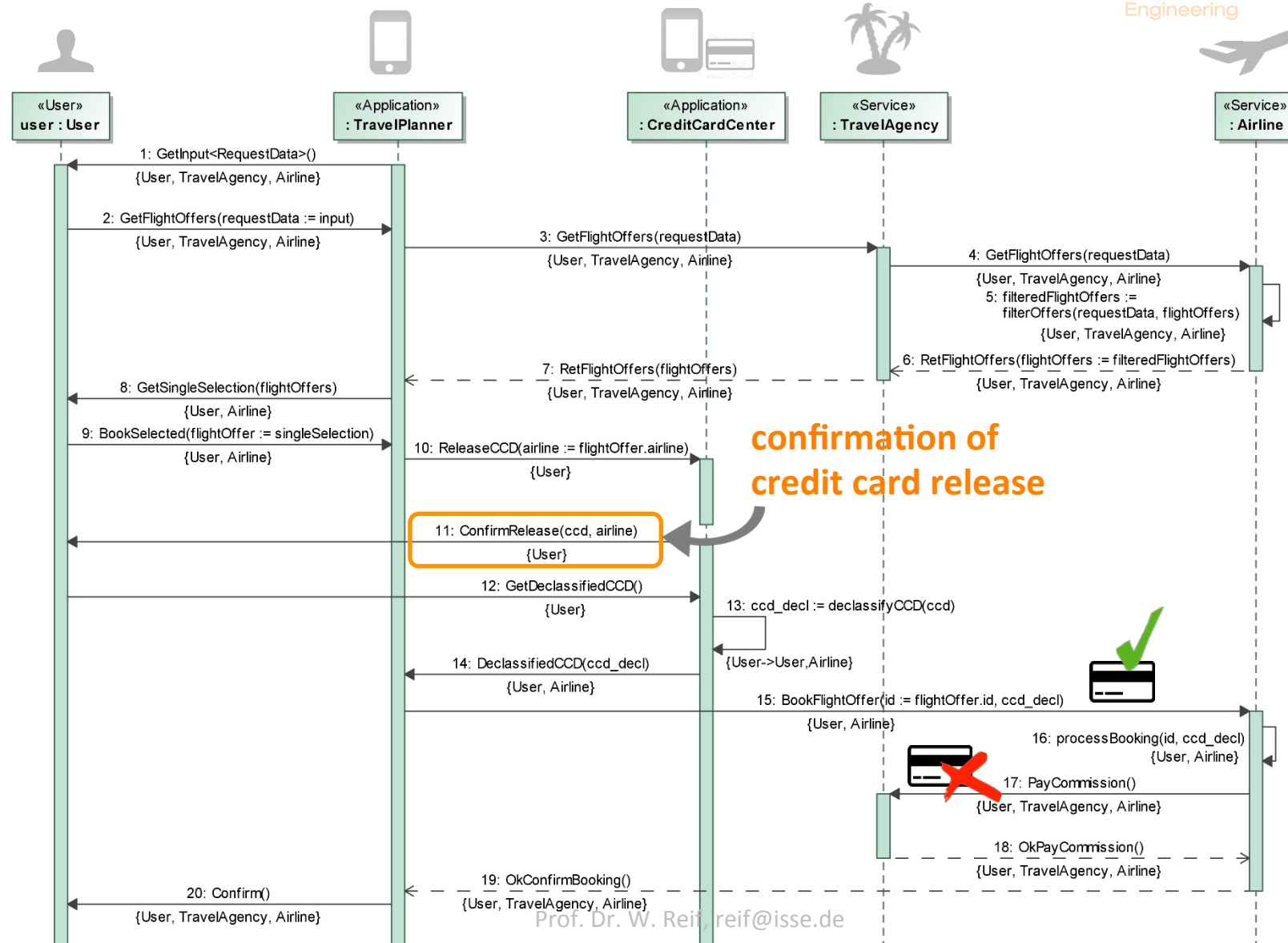


Travel Planner: security (cont.)

Assigning security domains to data & messages:



TravelPlanner: behaviour



Noninterference theorem

(extension of Rushby/van der Meyden):

Security policy satisfied, if

RM1: $s_1 \approx_d s_2 \rightarrow \text{output}(s_1, d) = \text{output}(s_2, d)$

(true by construction) ✓

RM2: $s_1 \approx_{\text{dom}(a)} s_2 \wedge s_1(l) = s_2(l) \wedge l \in \text{alter}(\text{dom}(a))$
 $\rightarrow \text{step}(s_1, a)(l) = \text{step}(s_2, a)(l)$

(true by construction) ✓

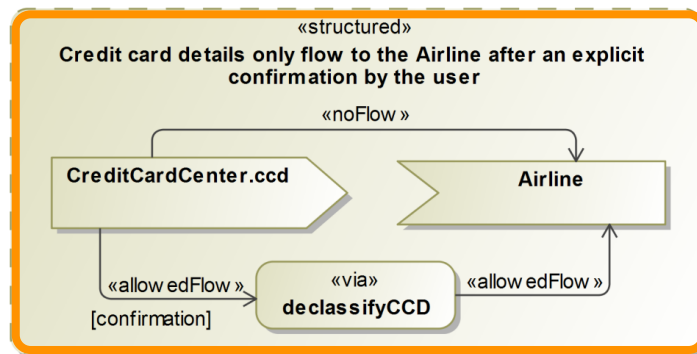
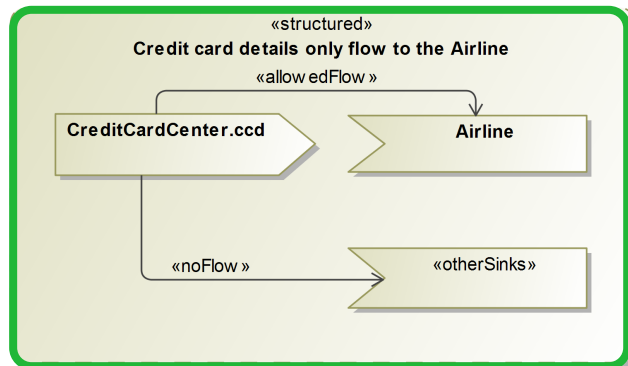
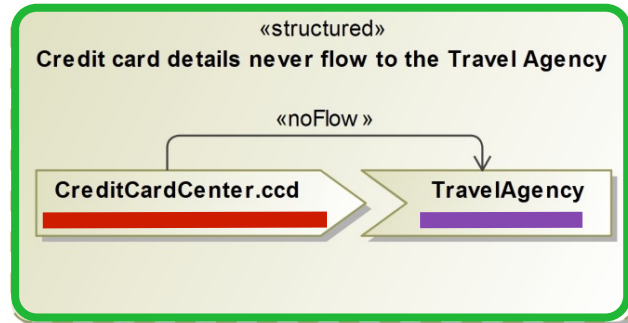
RM3: $\text{step}(s, a)(l) \neq s(l) \rightarrow l \in \text{alter}(\text{dom}(a))$

(true by construction) ✓

AOI: $\text{alter}(d_1) \cap \text{observe}(d_2) \neq \emptyset \rightarrow d_1 \rightsquigarrow d_2$

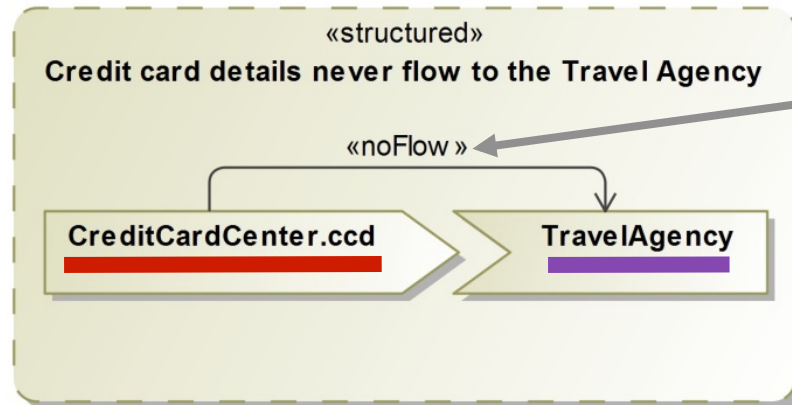
(prove)

TravelPlanner: security properties



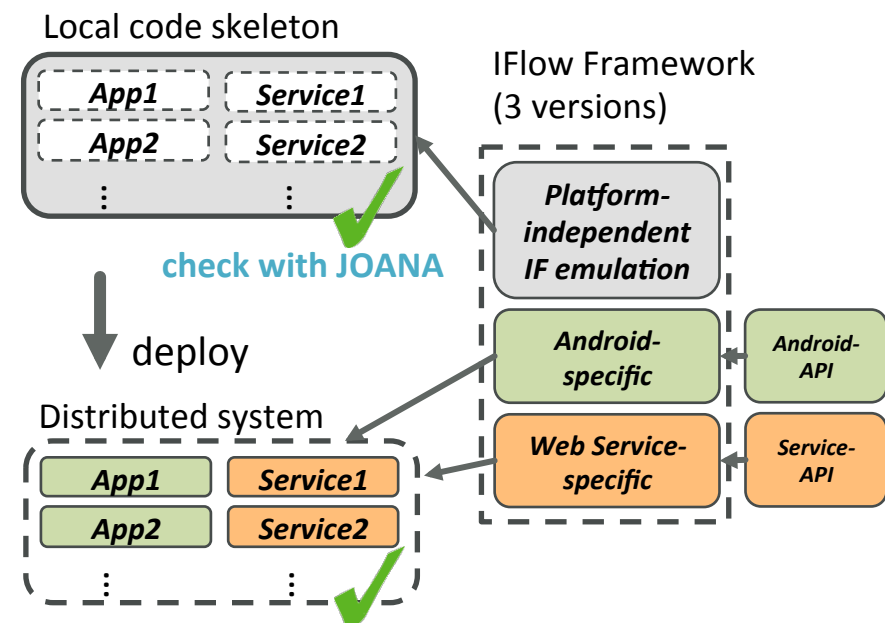
- Components (apps/services) and their attributes can be **sources** and **sinks** of information
- Information flow is either
 - **disallowed** («noFlow»), or
 - **allowed** («allowedFlow») ... via a specific method («via») ... after **user confirmation** ([confirmation])
- Properties can be
 - **automatically checked** on the code level, or
 - **verified interactively** using the formal model

TravelPlanner: automatic IF check

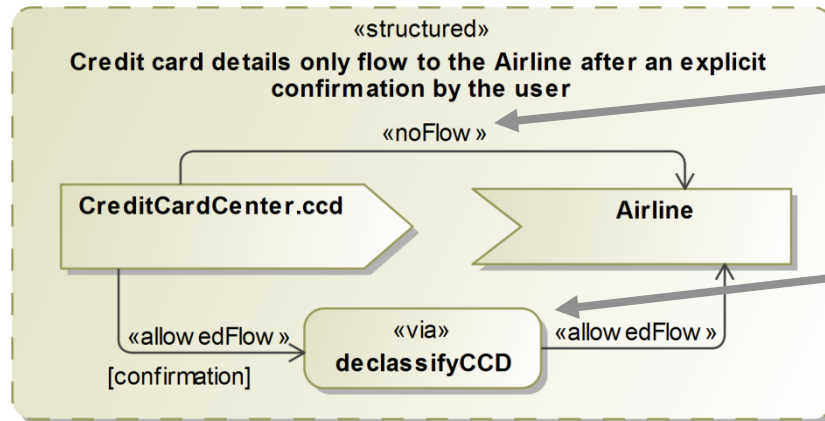


disallowed information flow

1. Generate a code skeleton as a **monolithic Java application**
2. Generate and run **information flow check** with JOANA (PDG-based IFC)
3. Distribute application as **Android apps** and **Java web services**



TravelPlanner: verification



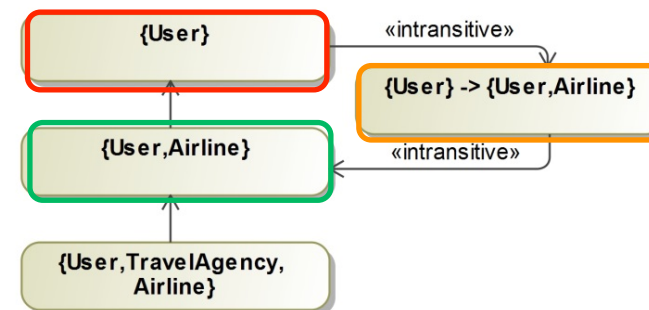
disallowed **direct** information flow

explicitly allowed information flow
via a specific method

Proof obligations:

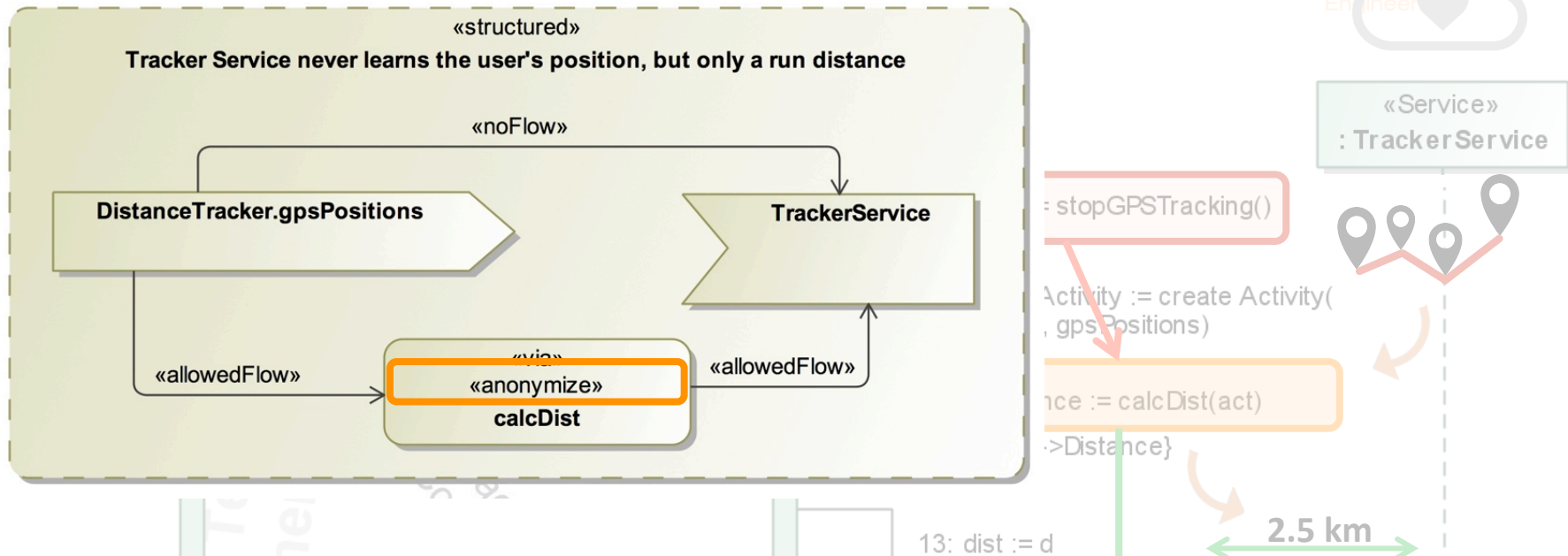
1. Security policy implies property
2. **Application model satisfies** security policy (noninterference theorem)
3. **declassifyCCD** is only called **after confirmation**

$\neg \text{domain}(\text{ccd}) \Rightarrow \text{domain}(\text{Airline}) \wedge$
 $\text{domain}(\text{ccd}) \Rightarrow \text{domain}(\text{declassifyCCD}) \wedge$
 $\text{domain}(\text{declassifyCCD}) \Rightarrow \text{domain}(\text{Airline})$



- IFlow: a model driven approach for developing distributed apps with secure information flow
- Software/security co-design in distributed apps
 - Modeling
 - Automatic check and verification
- **Special aspect: limited release of information**
 - Modeling
 - Verification

Distance Tracker: **property**



- Tracker Service **never learns** the user's location
- Distance Tracker **properly anonymizes** the user's location data as distance

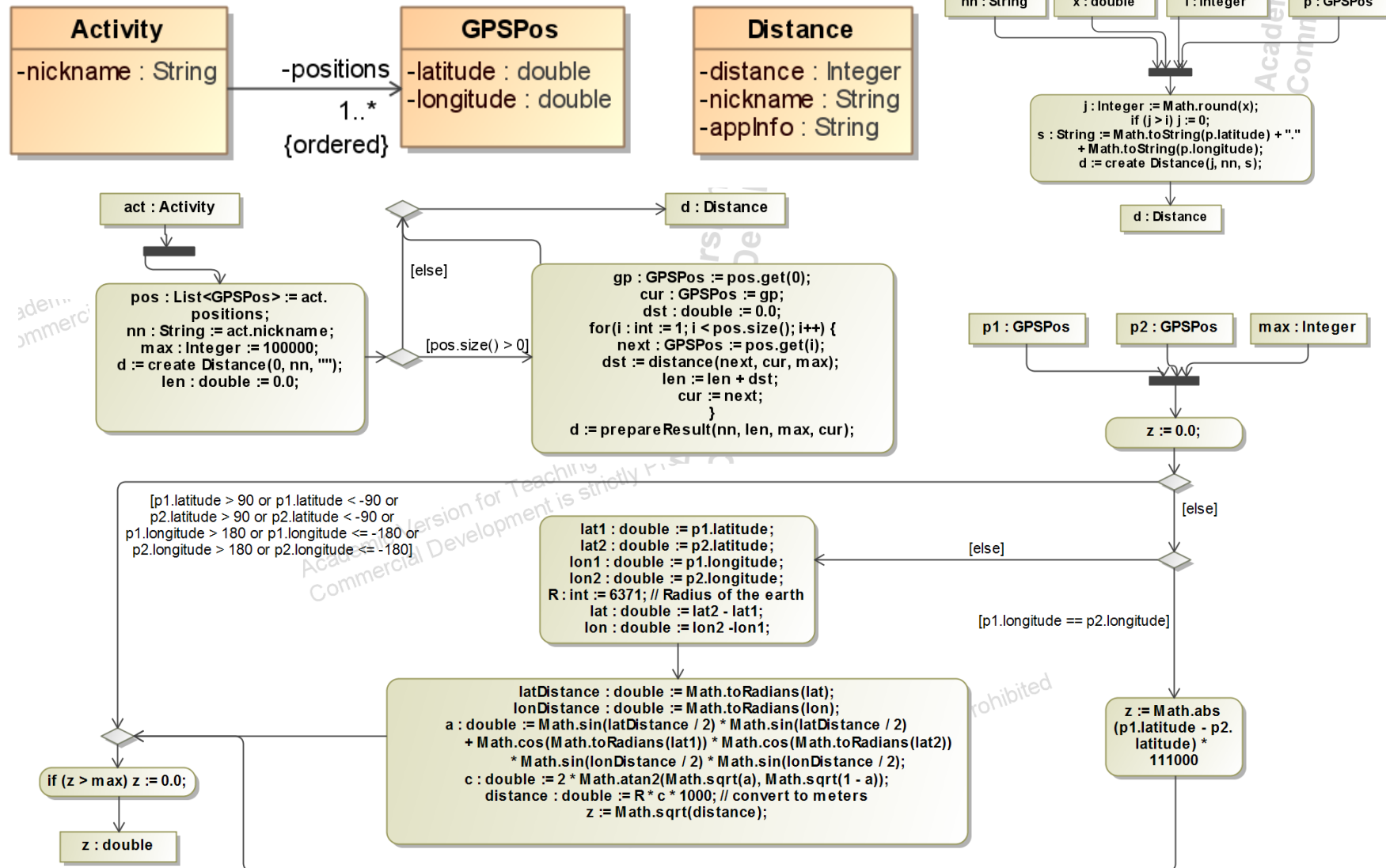
What do we prove?

- define equivalence class of all inputs that map to the same output
- reason about size of the classes

„For every valid output there are enough (***e***) different inputs with **mutually disjoint positions**“:

$$\begin{aligned} &\forall x, d. \text{calcDist}(x) = d \rightarrow \\ &\exists a_set. \# a_set \geq \mathbf{e} \wedge \mathbf{disjoint}(a_set) \wedge \\ &\forall y. y \in a_set \rightarrow \text{calcDist}(y) = d \end{aligned}$$

calcDist(act : Activity) : Distance



IFlow integrates

- a model driven approach (ModelFlow)
- formally verified information flow properties
- automatic information flow control